# **IDT**: **I**ntelligent **D**ata Placement for Multi-**t**iered Main Memory with Reinforcement Learning

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#### **Tiered Memory Systems**

- Emerging memory technologies are introducing **multiple tiers** in the **main memory**
	- ─ CXL Memory, HBM-enabled processors, Intel Optane DCPMM, …



#### **OS-level Tiered Memory Management**

• **OS** kernel manages **data placement** across tiers



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- **OS** kernel manages **data placement** across tiers
- **OS kernel demotes cold** pages to **lower**-tier memory



Main memory

Accurately **identifying data hotness** and effective **demotion criteria** are necessary!

#### **OS-level Tiered Memory Management**

- **OS** kernel manages **data placement** across tiers
- **OS kernel demotes cold** pages to **lower**-tier memory
- **OS kernel promotes hot** pages to **upper**-tier memory



#### Main memory

- Effective demotion candidate selection is crucial
	- ─ Impacts **promotion**
	- ─Incorrectly identifying demotion targets causes **ping-pong** of demotion and promotion
- Prior works used Linux kernel's **active/inactive LRU lists** (**2Q LRU)**
	- ─ Since 2022, **multi-generational LRU lists[1]** (**MGLRU)** for more fine-grained policy

• However, **2Q LRU** and **MGLRU** often deviate from the **actual data hotness** (PTE access bit scanning)



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**Cumulative probability distribution of accessed page's age varies across workloads**



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[1] Thaleia Dimitra Doudali et al., "Kleio: A Hybrid Memory Page Scheduler with Machine Intelligence," HPDC, 2019 [2] Andres Lagar-Cavilla et al., "Software-Defined Far Memory in Warehouse-Scale Computers," ASPLOS. 2019



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**IDT: Design and Implementation** 

#### **IDT: Overview**



#### IDT Overview



#### **Memory Access Monitoring**

#### **Region-granularity Monitoring**

- Monitor **group** of pages with **similar access patterns**
	- ─ **Partition** Virtual Memory Area (VMA) into **regions**



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- **Sample 2 pages** at each **sample\_interval**



#### **Region-granularity Monitoring**

- Monitor **group** of pages with **similar access patterns**
	- ─ **Partition** Virtual Memory Area (VMA) into **regions**
- **Sample 2 pages** at each **sample\_interval**
	- ─ Manage **history**, **access**, **age**[1]



#### **Region Reconfiguration**

- **Merge** or **split** adjacent regions for reconfiguration at each aggregate\_interval
	- ─ **Merge** regions with **similar access patterns** to **reduce** monitoring overhead



#### **Region Reconfiguration**

- **Merge** or **split** adjacent regions for reconfiguration at each aggregate\_interval
	- ─ **Merge** regions with **similar access patterns** to **reduce** monitoring overhead
	- ─ **Split** when pages in a region have **different access patterns**



#### **Region Reconfiguration**

• **Merge** or **split** adjacent regions for reconfiguration at each aggregate\_interval



#### **Region Reconfiguration: Merge**

• Validate the similarity of region's history vector by **Fisher's exact test** with a 90% significance level





window  $size = n$ 

#### **Region Reconfiguration: Merge**

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- **Sliding window** → Compare the **access ratio** of each region's window



#### **Region Reconfiguration: Merge**

- Validate the similarity of region's history vector by **Fisher's exact test** with a 90% significance level
- **Sliding window** → Compare the **access ratio** of each region's window
	- ─ If every window yields a *similar access ratio* → **Merge**





window  $size = n$ 



#### **Region Reconfiguration: Split**

• **Split** region when the **access status of the sampling pages differs** at sample\_interval



#### IDT Overview



#### **RL-based Demotion Policy**

#### **RL: Recall**



### **RL: Design**



### **RL: Design**



 $r(t) = log(demoted_pages(t) / slow_hitt(t))$ 

### **RL: Detail**

- Input Layer
	- $-$  min, q1 (25 percentile), q2 (50 percentile), q3 (75 percentile), max age distribution
	- ─ 1x5 state vector
- **2 Hidden Layers**
	- $-16$ , 32 nodes
- **Proximal Policy Optimization[1] (PPO)** Training Algorithm
- Experience buffer size: 4
	- ─ Trained every 4 inferences
- **Pre-train** with GUPS microbenchmark
	- $-$  3 memory access patterns used in HeMem<sup>[2]</sup>
- Implemented with PyTorch-based Rllib



#### **RL: Example**



#### **RL: Execution Phases**



#### **RL Execution Phases**



#### IDT Overview



#### **Demotion, Promotion Mechanism**

#### **Demotion**

Slower

- When a memory node's available space < demote\_wmark (Set to 10%)
- Demote regions with age > age\_thres and minimum access



#### **Promotion**

- **ARP** (**Accessed R**egion **P**romotion)
- **PRP** (**Predictive R**egion **P**romotion)

#### **Promotion: ARP (Accessed Region Promotion)**

- Promote when demoted region is **accessed**
	- ─ Destination node should have available space > critical\_wmark (Set to 1%)



#### **Promotion: PRP (Predictive Region Promotion)**

- ARP does not promote until access to the region's sampling pages is observed
	- ─ Preemptively promoting regions similar to ARP region may be beneficial
- Identify a **similar** region with **k-Nearest Neighbor** and promote

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*distance* **= Normalized(vaddr distance) + Normalized(access\_history distance) Spatial Locality Community Communist Spatial Locality** 

#### **More Details in the Paper**

- Aggressive demotion
	- ─ Tighten demotion criteria when scarce fast memory
- Misplaced region promotion
	- ─ Handle promotion of regions demoted by kswapd
- RL formulation
	- ─ Problem formulation
	- ─ Approximation for feasible implementation
- Sensitivity study

#### IDT: Intelligent Data Placement for Multi-tiered Main Memory with Reinforcement Learning



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#### **ABSTRACT**

To address the limitation of a DRAM-based single-tier in satisfying the comprehensive demands of main memory, multi-tiered memory systems are gaining widespread adoption. To support these systems, operating-system-level solutions that analyze the application's memory access patterns and ensure data placement in the appropriate memory tier have been vastly explored.

In this paper, we identify reinforcement learning (RL) as an effect tive solution for tiered memory management, and its policy can be formulated in a solvable form using RL. We also demonstrate that an effective region-granularity memory access monitoring method is necessary to provide an accurate environment state to the RL model. Thus, we propose IDT, an intelligent data placement for multitiered main memory. IDT incorporates an RL-based demotion policy autotuning and a mechanism that efficiently demotes cold pages to lower-tier memory. IDT also promotes hot pages to upper-tier memory to minimize access on slow memory, featuring a lightweight machine learning algorithm. IDT employs region-granularity memory access monitoring with statistical-testing-based adjacent region merge and split to improve precision and mitigate ambiguity observed in prior works. Experiments on an actual four-tiered memory system show that IDT achieves an average 2.08× speedup over the default Linux kernel and 11.2% performance improvement compared to the state-of-the-art solution.

#### **CCS CONCEPTS**

 $\cdot$  Software and its engineering  $\rightarrow$  Memory management; Computer systems organization  $\rightarrow$  Heterogeneous (hybrid)  $\textit{systems}\text{;}\textit{-}\textit{Computing methodologies}\rightarrow\textit{Reinforcement learn-}$ ing.

#### **KEYWORDS**

Memory Tiering, Emerging Memory Technologies, Memory Management, Reinforcement Learning

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#### 1 INTRODUCTION

The growing demand for memory-intensive workloads, such as high-performance computing, graph analytics, and in-memory databases, is highlighting the scaling limitations of a DRAM-based single-tier main memory [39]. To tackle this issue, a variety of memory types with diverse performance characteristics have been adopted to compose tiered memory systems. Recently, the rising interest in memories attached to Compute Express Link (CXL-Memory [9]) underscores that the future lies in multi-tiered memory systems by integrating various heterogeneous memories with a main-memory-like appearance to a system [36]. Cloud vendors, such as Amazon and Google, already serve large memory cloud instances based on multi-tiered memory systems [20, 33].

Tiered memory management requires a keen insight into an application's memory usage and placing the data in the proper memory tier according to its hotness. Thus, a number of prior studies have proposed operating-system-(OS)-level solutions to improve application performance by attentively exploiting the tiered mem ory system [2, 12, 15, 19, 23, 27, 36, 38, 48, 51, 55]. These OS-level tiered memory solutions typically consist of *data placement* to fully leverage diverse memory types and memory access monitoring to gather information for guiding data placement.

Data placement. Infrequently accessed pages in tiered memory should be *demoted* to lower-tier slow memory for efficient utiliza tion of upper-tier fast memory. Moreover, to complement demotion, hot pages trapped in slow memory should be identified and promoted to upper-tier memory. Several prior studies have utilized the Linux kernel's 2Q LRU [19, 21, 35, 36, 56, 57] or multi-generational LRU (MGLRU) [58] to determine demotion candidates. However, the data hotness identified by 2Q LRU or MGLRU often fails to reflect the actual data hotness of the application. Therefore, precisely tracking both access frequency and recency for each page, and establishing a demotion policy with solid criteria would be more effective. Yet, implementing this method presents the challenge of

#### **Evaluation**

#### **Experimental Setup**

- Based on Linux kernel v6.0.19
	- ─ Memory access monitoring developed with DAMON
- **Multi-tiered memory** setup
	- ─ 2 socket machine with **DRAM** (fast memory) and Intel Optane **DCPMM** (slow memory)
- 4 State-of-the-art solutions for comparison
	- ─ **Intel Tiering 0.8[1], TPP[2] , AutoTiering[3] , AutoNUMA Tiering (MGRLU)[4]**
- Workloads: SPEC CPU2017, graph500, GAPBS ─ RSS set **96GB~110GB** to facilitate using 3 tiers
- Evaluation metric: Speedup (execution time) normalized against AutoNUMA Balancing



[1] Intel. 2022. Tiering-0.8. [https://git.kernel.org/pub/scm/linux/kernel/git/vishal/tiering.git/.](https://git.kernel.org/pub/scm/linux/kernel/git/vishal/tiering.git/)

[2] Hasan Al Maruf et al., "TPP: Transparent Page Placement for CXL-Enabled Tiered-Memory," ASPLOS, 2023

[3] Jonghyeon Kim et al., "Exploring the Design Space of Page Management for Multi-Tiered Memory Systems," USENIC ATC (Virtual Event), 2021

[4] Yu Zhao. 2022. Multigenerational LRU Framework. [https://lwn.net/Articles/880393/.](https://lwn.net/Articles/880393/)

#### **Performance**

- Outperforms the best-performing state-of-the-art solution AutoNUMA Tiering (MGLRU)) by **11.2%**
	- ─ Average **2.08x** speedup against AutoNUMA Balancing



#### **Limitations**

- Other parameters (e.g. 10% and 1% watermarks, sliding window size) are not determined by RL (or ML)
	- ─ Our **goal was to advance the state-of-the-art** solution by **appropriately utilizing RL** (or ML)
	- ─ Future works may apply ML to optimize other parameters
- **Blackbox**: Difficult to explain clear reasons for performance improvement by using RL

### **Summary**

Inaccurate data hotness determined from 2Q/MGLRU

**ML** for selecting demotion candidates

**Statistical testing** for region-granularity monitoring

### **Summary**



Outperforms the default Linux kernel by **2.08×**, state-of-the-art solution by **11.2%**

# **Thank you!**

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**Backup Slides**

#### **RL Effectiveness**

- RL outperforms against static age\_thres
	- ─ When setting age\_thres to q2, q3, max of age distribution (Potential RL actions)



#### **RL Effectiveness (cont'd)**

- Performance variation on hyperparameters
	- $-$  Learning rate ( $\alpha$ ): Improvement over  $\alpha$ =0 shows efficacy of online training
	- $-$  **Exploration rate (** $\epsilon$ ): Improvement over  $\epsilon$ =1 shows effective than random policy
	- $\blacksquare$  **Discount factor (y)**: Improvement over  $\gamma=0$  shows effective than only accounting immediate reward



#### **Memory Access Monitoring Effectiveness**

- Compare against applying DAMON<sup>[1]</sup>
	- ─ Average history vector's hamming distance of merge region: 8.13 (DAMON) → **5.15 (IDT)**



**ZIDT-DAMON OIDT** 

#### **Predictive Region Promotion Effectiveness**

• Compare against without PRP



#### **Predictive Region Promotion Effectiveness (cont'd)**

- Compare against without PRP
	- ─ High PRP accuracy (ratio of region accessed that was promoted by PRP)



#### **Two-tiered memory**

- Performance on two-tiered memory configuration
	- ─ Set DRAM 0 (Tier 0) and DRAM 1 (Tier 1) to 64GB and RSS to 96GB~110GB
	- ─ Similar performance to AutoNUMA Tiering (MGRLU)



#### **Sensitivity Study**

- Smaller interval: Finer sampling and responsive demotion/promotion
	- ─ **but** overhead increase
- Larger watermark: Reserve fast memory for potential allocation requests
	- ─ **but** may not fully leverage the performance benefits of fast memory





#### **RL Pretraining**

- Pre-trained using the Giga Update operations Per Second (GUPS) microbenchmark<sup>[1]</sup> with 100GB RSS
	- ─ **Uniform random access**: Random access over the working set
	- ─ **Hot set**: 90% of access on 4GB hot objects and the remaining uniform randomly
	- ─ **Dynamic hot set**: Change hot objects every 150-second intervals.

#### **Aggressive Demotion**

- If available space < critical\_wmark (Set to 1%)
- **Tighten demotion candidate criteria**
	- $-$  age  $>$  age thres
	- ─ **access < (min\_access + max\_access) / 2**



#### **Misplaced Region Promotion**

- IDT's promotion may place region in suboptimal tier
- kswapd may demote in intensive memory usage
- Track by setting demoted flag when region is demoted
	- 1. Detect misplaced region with demoted flag
	- 2. Check upper-tier available space > critical\_wmark
	- 3. Promote

#### **Region Reconfiguration Methods**

- DAMON<sup>[1]</sup>
	- Merge if access frequency difference less than 10% of the maximum frequency across all regions
	- ─ Split randomly
- $MTM^{[2]}$ 
	- Merge if access frequency difference less than 1/3 of total scan counts
	- ─ Split if two sampling page access status differ